

# ESc 101: FUNDAMENTALS OF COMPUTING

## Lecture 25

Mar 8, 2010

# OUTLINE

## 1 MATRIX OPERATIONS

# BASIC MATRIX OPERATIONS

- Addition, subtraction
- Multiplication
- Inversion
- Computing Determinant

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# BASIC MATRIX OPERATIONS

- **Addition, subtraction:** Simple
- Multiplication
- Inversion
- Computing Determinant

# BASIC MATRIX OPERATIONS

- Addition, subtraction
- **Multiplication**: Done
- Inversion
- Computing Determinant



# BASIC MATRIX OPERATIONS

- Addition, subtraction
- Multiplication
- **Inversion**: Will develop an algorithm
- Computing Determinant

# BASIC MATRIX OPERATIONS

- Addition, subtraction
- Multiplication
- Inversion
- **Computing Determinant**: By definition, by Gaussian elimination

# DETERMINANT

Let  $A = [a_{i,j}]$  be an  $n \times n$  matrix. Its **determinant** is:

$$\sum_{\pi} \text{sgn}(\pi) \cdot \prod_{i=0}^{n-1} a_{i,\pi(i)},$$

where

- $\pi$  runs over all **permutations** of  $\{0, 1, 2, \dots, n-1\}$ , and
- $\text{sgn}(\pi) \in \{1, -1\}$  is the **sign** of permutation  $\pi$ .

# COMPUTING DETERMINANT

- Computing determinant using the above formula will be very time consuming: as there are  $n!$  permutations of  $\{0, 1, 2, \dots, n - 1\}$ , and the formula sums over all of these.
- There is a faster way known for computing determinant: Gaussian elimination.

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- There is a faster way known for computing determinant: **Gaussian elimination**.

# GAUSSIAN ELIMINATION

Let  $A_0 = [a_{i,j}]$  be an  $n \times n$  matrix:

$$A_0 = \begin{bmatrix} a_{0,0} & a_{0,1} & \cdots & a_{0,n-1} \\ a_{1,0} & a_{1,1} & \cdots & a_{1,n-1} \\ \vdots & \vdots & \ddots & \vdots \\ a_{n-1,0} & a_{n-1,1} & \cdots & a_{n-1,n-1} \end{bmatrix}.$$

**FIRST STEP:** Check if  $a_{0,0} \neq 0$ . If it is, add to it the first row whose first element is non-zero. If no such row exists, then the determinant is **zero**.

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# GAUSSIAN ELIMINATION

**SECOND STEP:** For every  $i > 0$ , subtract  $\frac{a_{i,0}}{a_{0,0}}$  times the first row from the  $i$ th row. This makes  $a_{i,0} = 0$  for all  $i > 0$ .

After the first two steps, the matrix looks like:

$$A_0 = \begin{bmatrix} a_{0,0} & a_{0,1} & \cdots & a_{0,n-1} \\ 0 & a_{1,1} & \cdots & a_{1,n-1} \\ \vdots & \vdots & \ddots & \vdots \\ 0 & a_{n-1,1} & \cdots & a_{n-1,n-1} \end{bmatrix},$$

where the values of many elements has been modified from their original value.



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Let matrix  $A_1$  be:

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NEXT STEPS: Repeat the first two steps for  $A_1$  and all the submatrices  $A_2, \dots, A_{n-1}$  that arise.

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**NEXT STEP:** Let matrix  $B$  be defined by taking the first row of  $A_0$ , second row of  $A_1$ , ..., last row of  $A_{n-1}$ .

Matrix  $B$  looks like:

$$B = \begin{bmatrix} a_{0,0} & a_{0,1} & a_{0,2} & \dots & a_{0,n-1} \\ 0 & a_{1,1} & a_{1,2} & \dots & a_{1,n-1} \\ 0 & 0 & a_{2,2} & \dots & a_{2,n-1} \\ \vdots & \vdots & \ddots & \ddots & \vdots \\ 0 & 0 & 0 & \dots & a_{n-1,n-1} \end{bmatrix}.$$

**LAST STEP:** The determinant of the matrix  $A$  equals the product of diagonals of  $B$ , i.e.,  $\prod_{i=0}^{n-1} a_{i,i}$ .

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# WHY DOES IT WORK?

## THEOREM

*The determinant of a matrix does not change by adding or subtracting a row to another row.*

The Gaussian Elimination algorithm only adds or subtracts rows.

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# REWRITING THE ALGORITHM

Input: matrix  $A$ , and its size  $n$ .

1. If  $(n == 1)$  go to step 5;
2. If  $(A[0][0] == 0)$  {  
    Find the smallest  $i$  such that  $A[i][0] != 0$ ;  
    If there is no such  $i$  then // determinant is zero  
        return 0;  
    Add row  $A[i]$  to row  $A[0]$ ;  
}
3. For every  $i > 0$ :  
    Replace row  $A[i]$  by  $A[i] - (A[i][0]/A[0][0]) * A[0]$ ;
4. Drop first row and first column of  $A$  and go back to 1;
5. Return the product of diagonal elements;

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## CONVERTING TO A C PROGRAM

```
/* Computes the determinant of size n matrix
 * stored in array A.
 */
float determinant(float A[][N], int n)
{
    float B[N][N]; // stores a submatrix of A
    int m; // the size of B
    float det = 1.0; // determinant value
    int i;

    copy_matrix(B, A, 0, n); // copy A to B
```

## CONVERTING TO A C PROGRAM

```
/* Do the Gaussian elimination for the first row,  
 * multiply the first diagonal element to det, and drop  
 * the first row and column from B.  
*/  
for (m = n; m > 0; m--) {  
    if (B[0][0] == 0) {  
        i = find_nonzero_row(B, m);  
        if (i >= m) // no non-zero row  
            return 0.0; // determinant is 0  
        add_row(B[0], B[i], 1, m); // add row i to row 0  
    }  
}
```

## CONVERTING TO A C PROGRAM

```
// Make first column of B zero except the first row
for (int t = 1; t < m; t++)
    add_row(B[t], B[0], - B[t][0]/B[0][0], m);

det = det * B[0][0]; // update determinant value

// drop the first row and column of B
copy_matrix(B, B, 1, m-1);
}
return det;
}
```



## CONVERTING TO A C PROGRAM

```
/* Copies matrix A to B after dropping first i rows and
 * columns of A. The size of matrix B is m.
 */
void copy_matrix(float B[][N], float A[][N], int i, int m)
{
    for (int k = 0; k < m; k++)
        for (int j = 0; j < m; j++)
            B[k][j] = A[k+i][j+i];
}
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# RECALL ALGORITHM

- Observe that after Step 4, we get a matrix of size  $n - 1$  and we need to compute its determinant.
- This is the same problem as the original one, except that the size is one less.
- So we can use the **same algorithm** to solve it.
- That is why, the execution goes back to Step 1.
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 */
float determinant(float A[][N], int n)
{
    float B[N][N]; // stores a submatrix of A
    int i;

    if (n == 1) // Step 1: 1 x 1 matrix
        return A[0][0];
}
```

# PROGRAM USING RECURSION

```
// Step 2: Make A[0][0] non-zero
if (A[0][0] == 0) {
    i = find_nonzero_row(A, n);
    if (i >= n) // no non-zero row
        return 0.0; // determinant is 0
    add_row(A[0], A[i], 1, n); // add row i to row 0
}
```

# PROGRAM USING RECURSION

```
// Step 3: Make first column of A zero except first row
for (int t = 1; t < n; t++)
    add_row(A[t], A[0], - A[t][0]/A[0][0], n);

// Step 4: drop the first row and column of A
copy_matrix(B, A, 1, n-1);

return A[0][0] * determinant(B, n-1); // recursive call!
}
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# PROGRAM USING RECURSION

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// Step 3: Make first column of A zero except first row
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return A[0][0] * determinant(B, n-1); // recursive call!
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# RECURSION

- A function is **recursive** if it is called inside its own definition.
- Such a definition is a substitute for loop, as in the example above.
- The execution jumps to the beginning of the function at the recursive call.
- To avoid infinite repetitions, it is necessary that:
  - ▶ in every successive call, some parameter value reduces,
  - ▶ and for small enough value of that parameter, there is no recursive call in the function.

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